



CALIFORNIA STATE UNIVERSITY  
**FULLERTON**

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# CPSC-440 Computer System Architecture

## Lecture 3

### Von Neumann Machines (IAS)

# History of Computers

## First Generation: Vacuum Tubes

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- ENIAC
  - Electronic Numerical Integrator And Computer
- Designed and constructed at the University of Pennsylvania
  - Started in 1943 – completed in 1946
  - By John Mauchly and John Eckert
- World's first general purpose electronic digital computer
  - Army's Ballistics Research Laboratory (BRL) needed a way to supply trajectory tables for new weapons accurately and within a reasonable time frame
  - Was not finished in time to be used in the war effort
- Its first task was to perform a series of calculations that were used to help determine the feasibility of the hydrogen bomb
- Continued to operate under BRL management until 1955 when it was disassembled



# ENIAC

Weighed  
30  
tons

Occupied  
1500  
square  
feet  
of  
floor  
space

Contained  
more  
than  
18,000  
vacuum  
tubes

140 kW  
Power  
consumption

Capable  
of  
5000  
additions  
per  
second

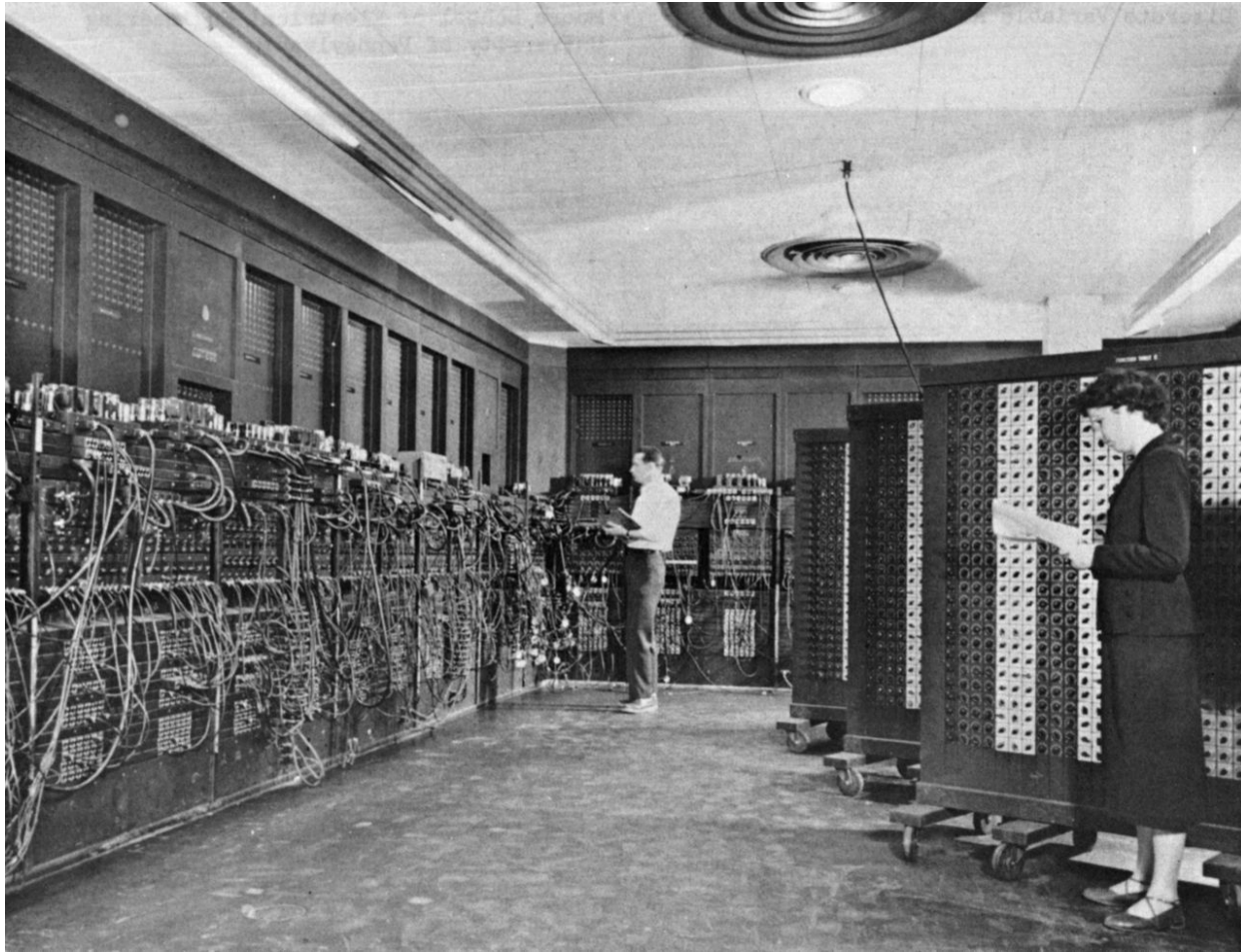
Decimal  
rather  
than  
binary  
machine

Memory  
consisted  
of 20  
accumulators,  
each  
capable  
of  
holding  
a  
10 digit  
number

Major  
drawback  
was the need  
for manual  
programming  
by setting  
switches  
and  
plugging/  
unplugging  
cables



# ENIAC



# John von Neumann

## EDVAC (Electronic Discrete Variable Computer)

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- First publication of the idea was in 1945
- Stored program concept
  - Attributed to ENIAC designers, most notably the mathematician John von Neumann
  - Program represented in a form suitable for storing in memory alongside the data
- IAS computer
  - Princeton Institute for Advanced Studies
  - Prototype of all subsequent general-purpose computers
  - Completed in 1952



# Structure of von Neumann Machine

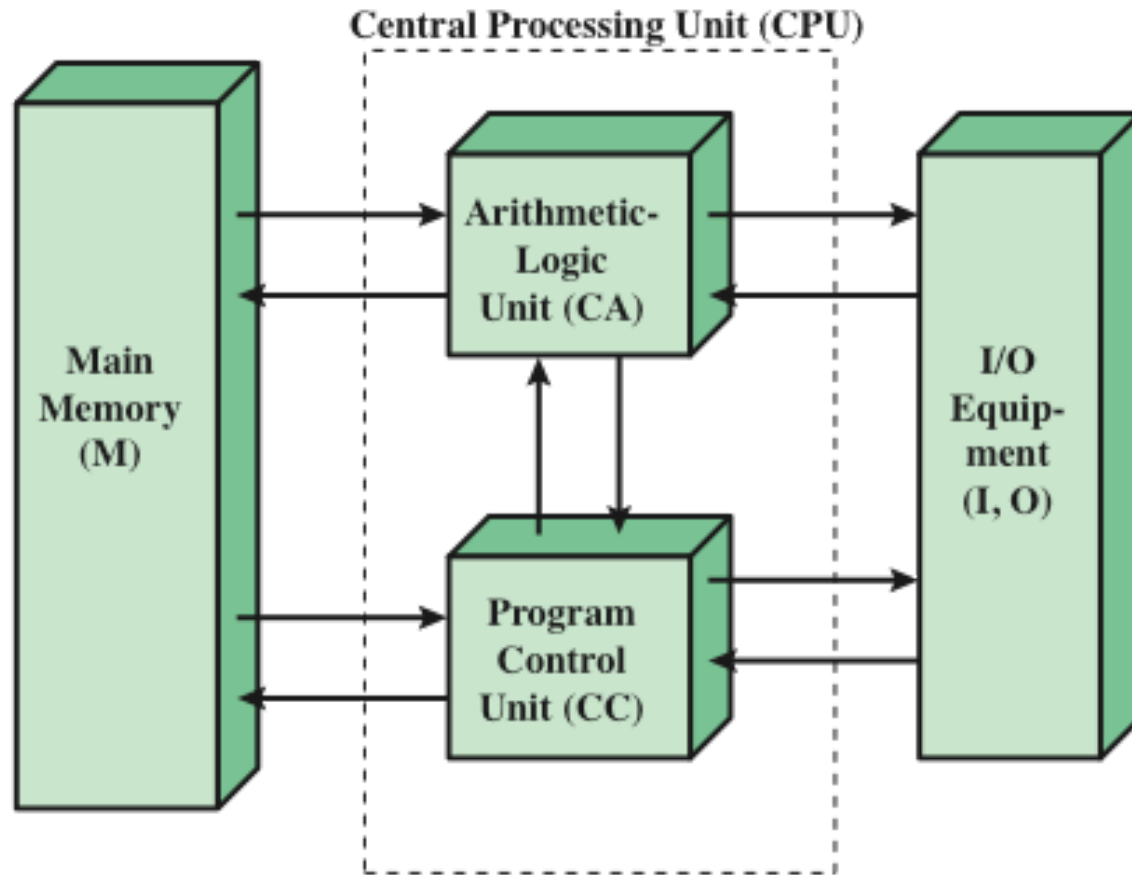


Figure 2.1 Structure of the IAS Computer



# IAS Memory Formats

- The memory of the IAS consists of 1000 storage locations (called words) of 40 bits each
- Both data and instructions are stored there
- Numbers are represented in binary form and each instruction is a binary code

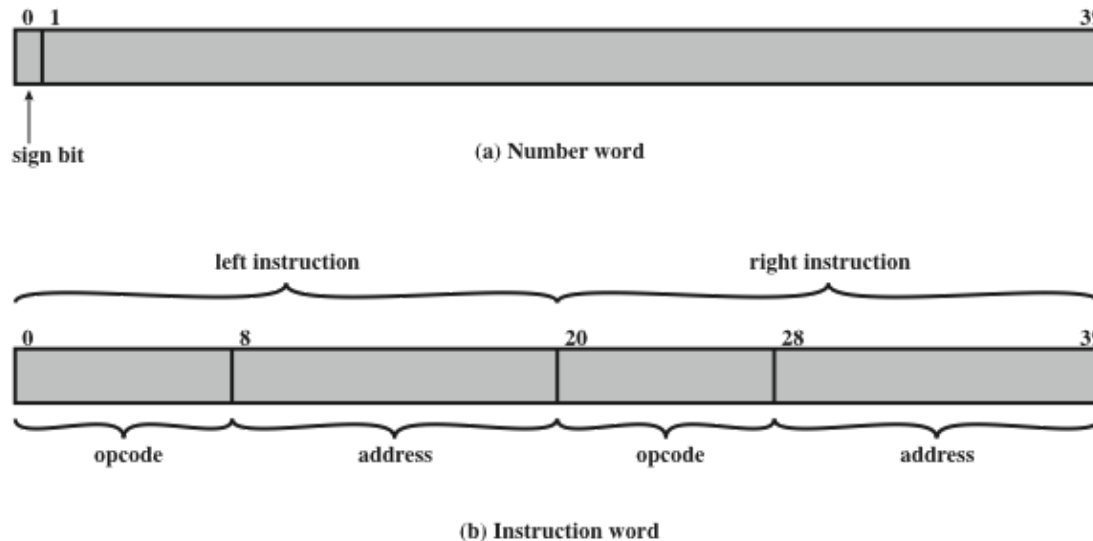


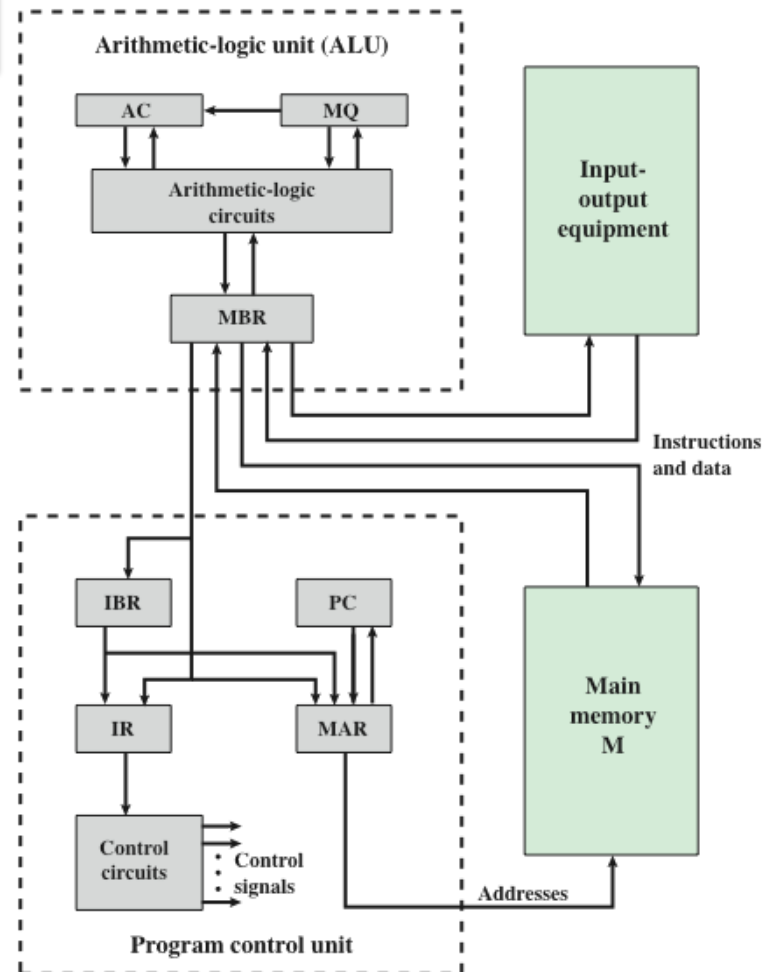
Figure 2.2 IAS Memory Formats



# Structure of IAS Computer

## Registers

- Memory Buffer Register (MBR)
  - Word to be stored/received in/from memory or I/O unit
- Memory Address Register (MAR)
  - Memory Address of the word to be (written from)/(read into) the MBR
- Instruction Register (IR)
  - Contains 8-bit opcode
- Instruction Buffer Register (IBR)
  - Temporarily holds the right-hand instruction
- Program Counter (PC)
  - Contains address of the next instruction pair to be fetched from memory
- Accumulator (AC) and Multiplier Quotient (MQ)
  - Employed to temporarily hold operands and results of ALU operations





Instruction Type	Opcode	Symbolic Representation	Description
Data transfer	00001010	LOAD MQ	Transfer contents of register MQ to the accumulator AC
	00001001	LOAD MQ,M(X)	Transfer contents of memory location X to MQ
	00100001	STOR M(X)	Transfer contents of accumulator to memory location X
	00000001	LOAD M(X)	Transfer M(X) to the accumulator
	00000010	LOAD -M(X)	Transfer -M(X) to the accumulator
	00000011	LOAD  M(X)	Transfer absolute value of M(X) to the accumulator
Unconditional branch	00000100	LOAD - M(X)	Transfer - M(X)  to the accumulator
	00001101	JUMP M(X,0:19)	Take next instruction from left half of M(X)
	00001110	JUMP M(X,20:39)	Take next instruction from right half of M(X)
Conditional branch	00001111	JUMP+ M(X,0:19)	If number in the accumulator is nonnegative, take next instruction from left half of M(X)
	00010000	JUMP+ M(X,20:39)	If number in the accumulator is nonnegative, take next instruction from right half of M(X)
Arithmetic	00000101	ADD M(X)	Add M(X) to AC; put the result in AC
	00000111	ADD  M(X)	Add  M(X)  to AC; put the result in AC
	00000110	SUB M(X)	Subtract M(X) from AC; put the result in AC
	00001000	SUB  M(X)	Subtract  M(X)  from AC; put the remainder in AC
	00001011	MUL M(X)	Multiply M(X) by MQ; put most significant bits of result in AC, put least significant bits in MQ
	00001100	DIV M(X)	Divide AC by M(X); put the quotient in MQ and the remainder in AC
	00010100	LSH	Multiply accumulator by 2; i.e., shift left one bit position
	00010101	RSH	Divide accumulator by 2; i.e., shift right one position
Address modify	00010010	STOR M(X,8:19)	Replace left address field at M(X) by 12 rightmost bits of AC
	00010011	STOR M(X,28:39)	Replace right address field at M(X) by 12 rightmost bits of AC

## Table 2.1

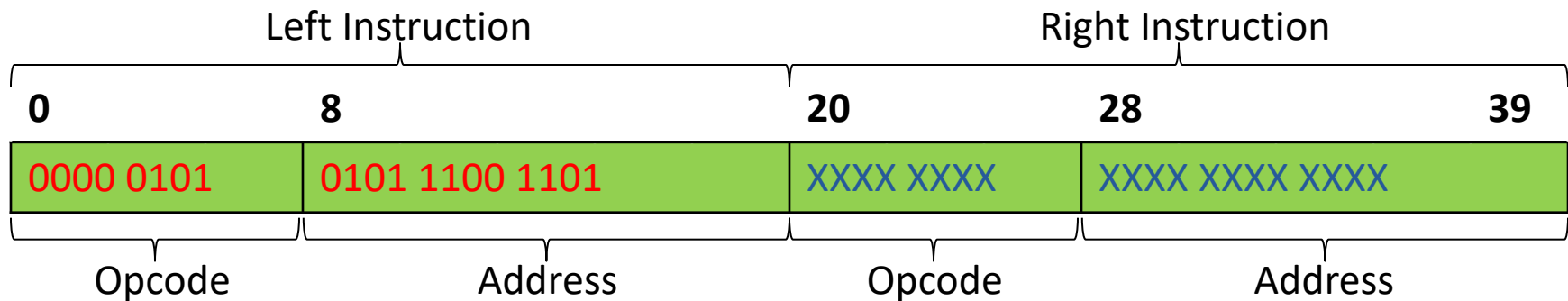
# The IAS Instruction Set

Table 2.1: The IAS Instruction Set



# Example 1

- What would the machine code instruction look like to add the contents of memory address 5CD (HEX) with the accumulator and stores the result back into the accumulator?

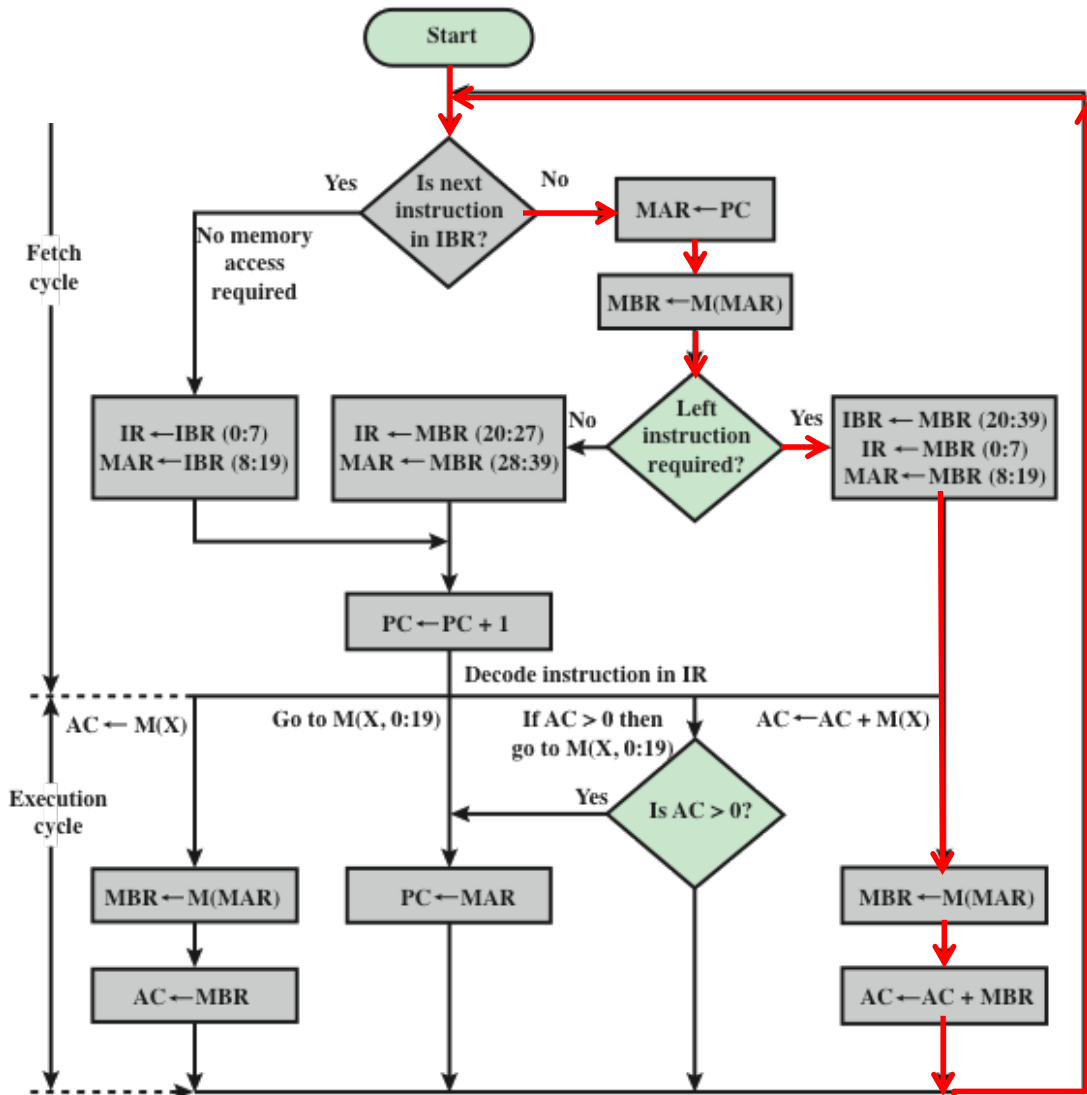


Symbolic	Description
ADD M(X)	Add M(X) to AC; put the result in AC



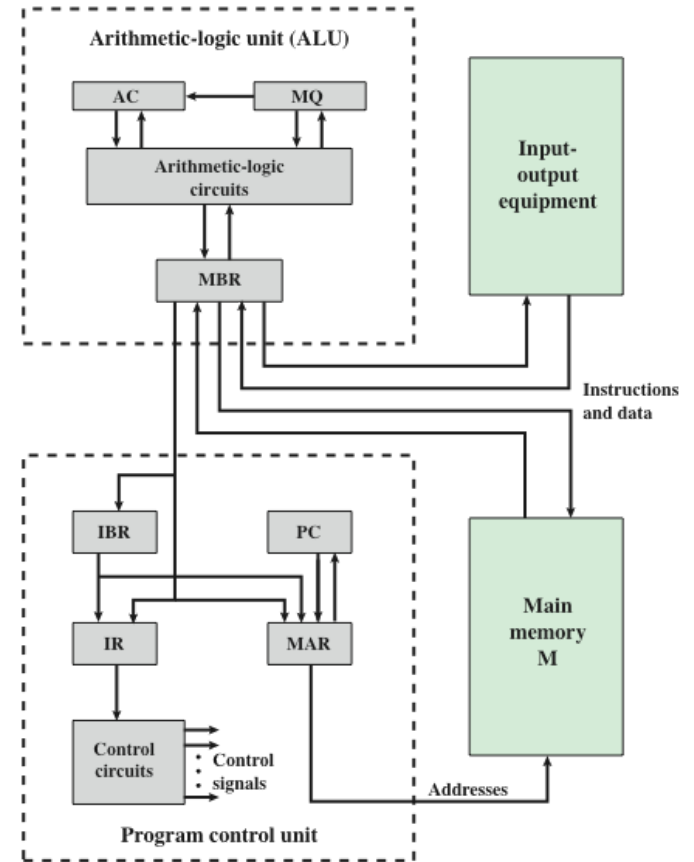
# Example 1

Left instruction first



M(X) = contents of memory location whose address is X  
(i:j) = bits i through j

Figure 2.4 Partial Flowchart of IAS Operation



# Example 1

The execution path would depend on the right instruction

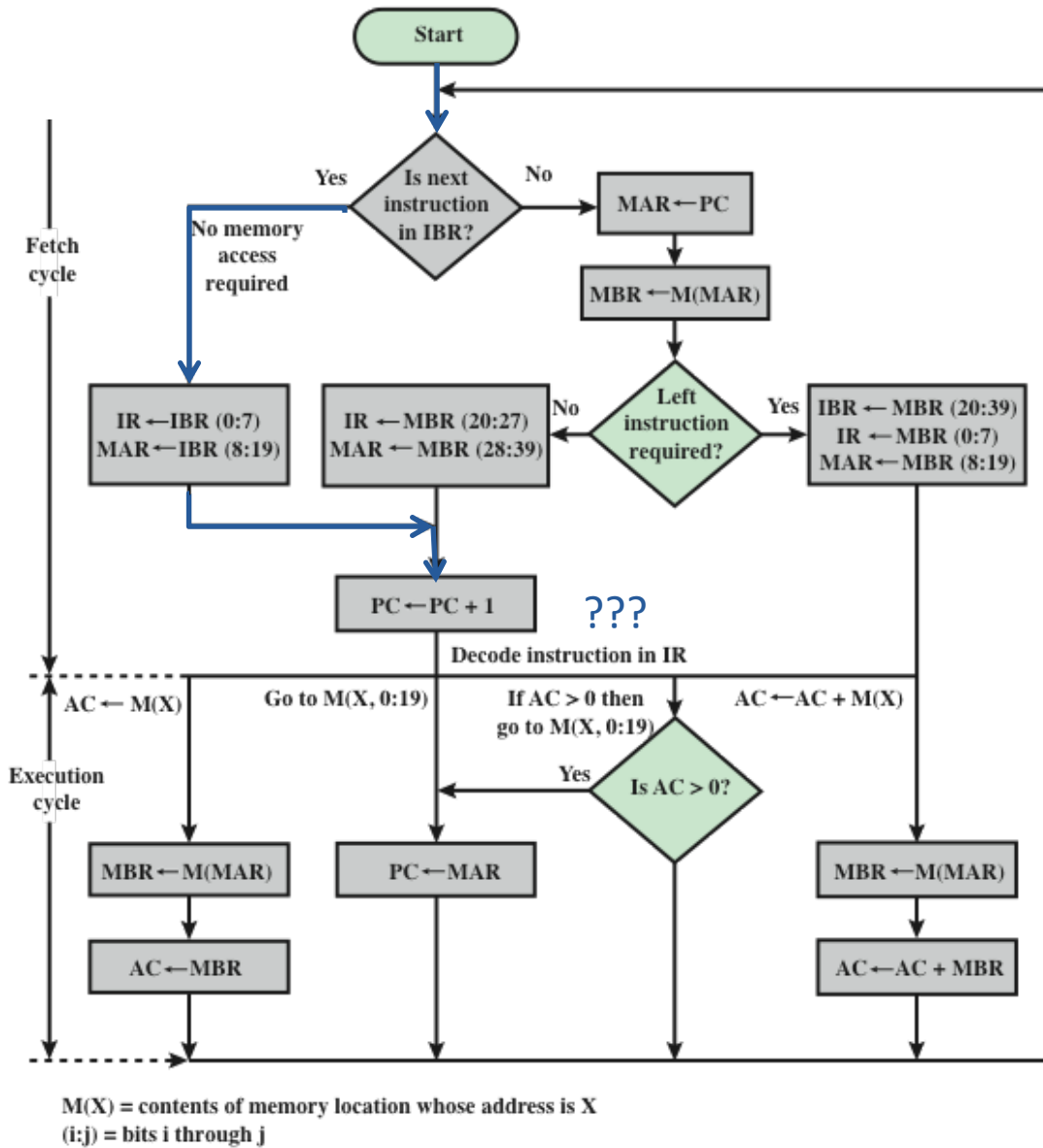
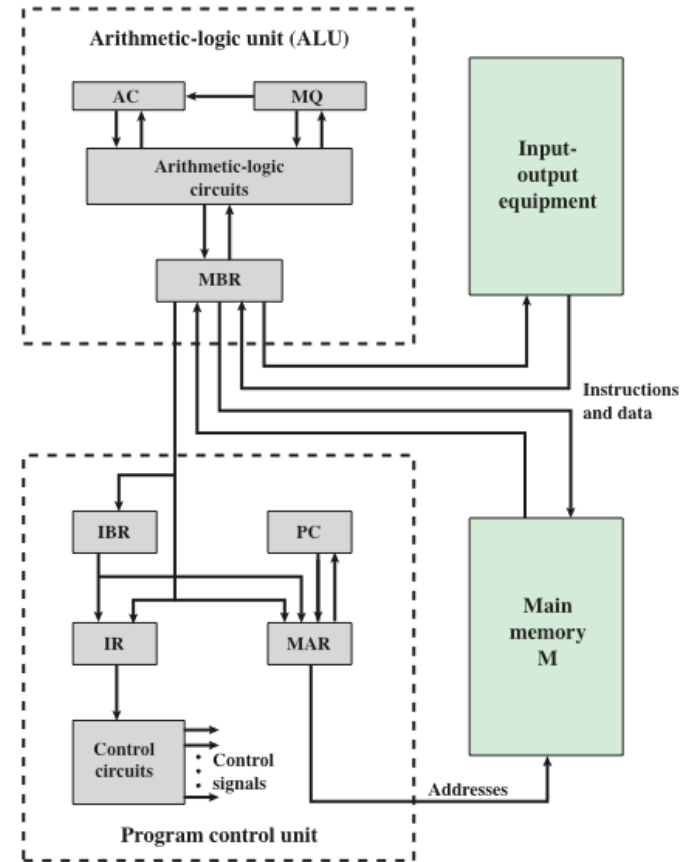
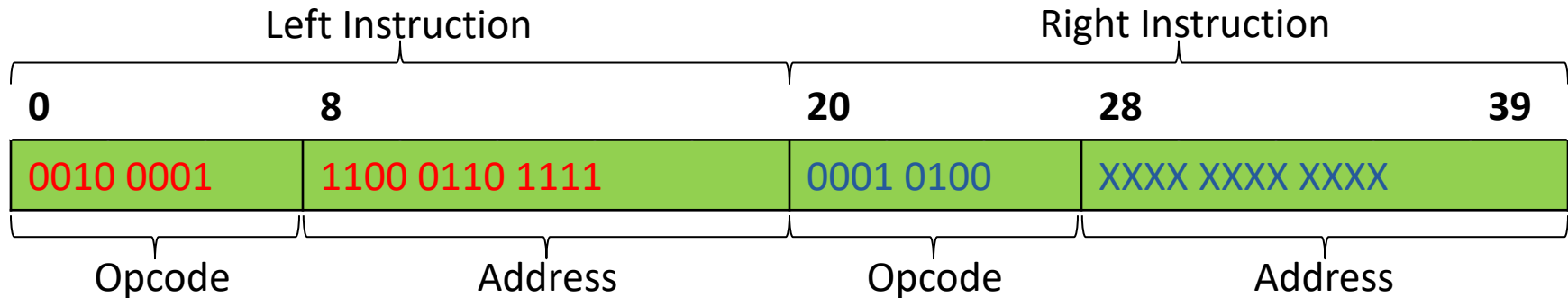


Figure 2.4 Partial Flowchart of IAS Operation



# Example 2

- What is the assembly language code for the program:

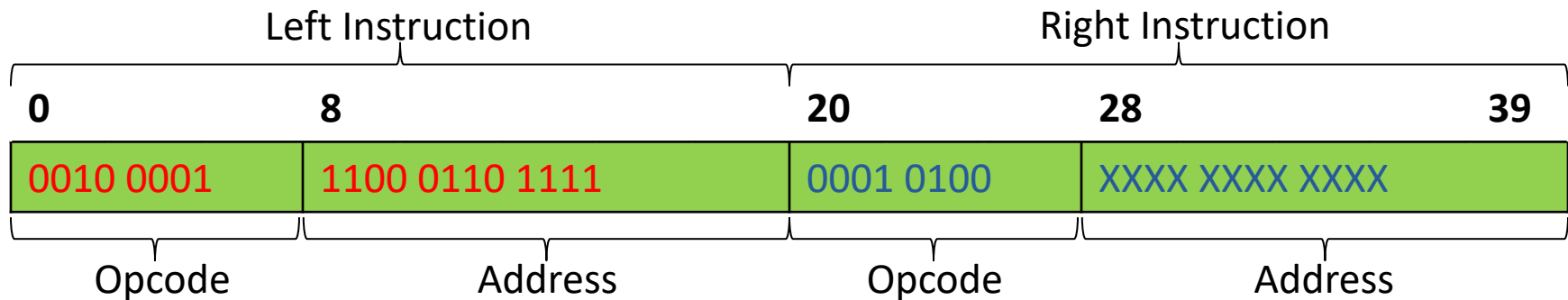


Address	Machine Code
06B	21C6F14XXX

NOTE: IAS doesn't actually have an assembly language

# Example 2

- What is the assembly language code for the program:



Address	Machine Code
06B	21C6F14XXX

Address	Symbolic
06B	STOR M(C6F)
	LSH

# Example 3

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- Write an IAS program to compute the results of the following equation:

$$Y = \frac{N(N + 1)}{2}$$

- Assume that the result of the computation doesn't overflow and N is a positive integer



# Example 3

$$Y = \frac{N(N + 1)}{2}$$

Location	Instruction/Value	Comments
0	<>	Constant (N) [initialized to some value]
1	1	Constant; Integer value = 1
2	2	Constant; Integer value = 2
3	0	Variable Y (initialized to integer zero)
4	0	Variable X (initialized to integer zero)
5L	LOAD M(0)	N → AC
5R	ADD M(1)	AC + 1 → AC; (N+1)
6L	STOR M(4)	AC → X; X=N+1
6R	LOAD MQ,M(4)	X → MQ; MQ=N+1
7L	MUL M(0)	MQ*M(0) → N(N+1) → AC
7R	DIV M(2)	AC/2 → AC; AC=N(N+1)/2
8L	STOR M(3)	AC → Y; saving the Sum in variable Y
8R	JUMP M(8,20:39)	Jump to 8R; Done





# Example 4

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- Write an IAS program to compute the results of the following equation:

$$Y = \sum_{X=1}^N X$$

- Assume that the result of the computation doesn't overflow, and that  $X$ ,  $Y$ , and  $N$  are positive integers



# Example 4

$$Y = \sum_{X=1}^N X$$

Location	Instruction/Value	Comments
0	<>	Constant (N) [initialized to some value]
1	1	Constant (loop counter increment)
2	1	Variable i (loop index value; current)
3	1	Variable Y = Sum of X values (Initialized to One)
4	LOAD M(0)	N → AC (the max limit)
5L	SUB M(2)	Compute N-i → AC
5R	JUMP + M(6,20:39)	If AC > 0 [i < N] then jump to 6R
6L	JUMP M(6,0:19)	Loop here (HALT)
6R	LOAD M(2)	i < N so continue; Get loop counter i
7L	ADD M(1)	i+1 in AC
7R	STOR M(2)	AC → i
8L	ADD M(3)	i + Y in AC
8R	STOR M(3)	AC → Y
9L	JUMP M(5,0:19)	Jump to 5L

# Homework Problems

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- Problems are available on Canvas