

# **Lecture Overview**



- Heuristic search algorithms
  - Hill-climbing,
  - Simulated annealing,
  - Best-first search and A\*,
  - Constraint satisfaction,
  - Mini-max for game playing
  - Advanced game playing
- Performance evaluation of heuristic search
  - Criteria for evaluation of heuristics
  - Complexity and efficiency issues
  - AI-complete and AI-hard

n	n!	
0	1	
1	1	
2	2	
3	6	
4	24	
5	120	
6	720	
7	5040	
8	40320	
9	362880	
10	3628800	
11	39916800	
12	479001600	
13	6227020800	
14	87178291200	
15	1307674368000	
16	20922789888000	
17	355687428096000	
18	6402373705728000	
19	121645100408832000	
20	2432902008176640000	
25	1.551121004×10 <sup>25</sup>	
50	3.041409320×10 <sup>64</sup>	
70	1.197857167×10 <sup>100</sup>	
100	9.332621544×10 <sup>157</sup>	
450	1.733368733×10 <sup>1000</sup>	
1000	4.023872601×10 <sup>2567</sup>	



Search space for Tic-Tac-Toe: 9!

Search space for chess: 64!

Search space for Go! game: 361!

 $10^{10} \text{ ms} = 115.7 \text{ days}$ 

# **Heuristics**



- Main problems of BFS and DFS
  - Brute force approach leading to combinatorial explosion
    - No utilization of information/strategies to make it efficient
- Many straightforward search algorithms can be improved by utilizing information and *intelligent strategy*, called "*heuristics*"
  - **Heuristics** is the study of the methods and rules of discovery and invention (Eureka "I have found it!"), commonly used by human.
  - Improvement can be made by intelligent strategies based on utilization of information or knowledge.
- Heuristic is necessary when there is no perfect solution.
  - E.g., playing chess or go game, many optimization problems, most AI problems.

# First Three Levels of the Tic-Tac-Toe State Space Reduced by Symmetry



# The "most wins" Heuristic Applied to the First Children in Tic-Tac-Toe







Three wins through a corner square

Four wins through the center square

Two wins through a side square

# Heuristically Reduced State Space for Tic-Tac-Toe



2/3 of all the search space is pruned away with the first move.

# Search Space for the Travelling Salesman Problem



# Looking for the Shortest Path with the Nearest Neighbor Path



Is it optimal solution?

\*Note: if (A,C) is 125, this path (A, E, D, B, C, A), at a cost of **375**, is the shortest path. If (A, C) is 300, then the cost is 550, not the shortest. The comparatively high cost of arc (C, A) defeated the heuristic (nearest neighbour) but it is better than most random paths.

# **How to Define Heuristics**



- General approaches to define heuristics
  - Identify the goal of the problem.
  - Collect and analyze available information and use background/domain knowledge to achieve the goal.
  - Identify all possible states and eliminate unnecessary states using the information to reduce the search space if possible.
  - Define a method(s) to quantitatively evaluate each state (heuristics).
  - Formulate the heuristics into a function (heuristic function).
- Try some examples
  - Tic-tac-toe, 8-puzzle, Chess, Go
  - Auto pilot for car, Mars rover, Robot vacuum

# **How to Define a Heuristic Function**



- A heuristic function, f(n) = g(n) + h(n)
  - **g(n)**: the measure (distance or cost) from the start to current state, n,
    - E.g., count 0 for the beginning state and is incremented by 1 for each level of the search.
  - h(n): a heuristic <u>estimate</u> of the measure from state n to a goal.
- h() guides search toward heuristically promising states while
   g() prevents search from indefinitely fruitless path.
  - If two states are the same or nearly the same, it is generally preferable to examine the state that is **nearest** to the root state of the graph (initial state) since it will give a greater probability of being on the shortest path to the goal.

### • Cost vs. Reward

- If **f()** is a cost function, the smaller the better.
- If f() is a reward function, the larger the better.



1	2	3
8		4
7	6	5
Goal		

#### Possible heuristics h():

h1() Counts the tiles out of place compared with the goal state in each state.
h2() Sum all the distances (Manhattan distance) by which the tiles are out of place.
h3() Multiplies a small number, e.g., 2 times each direct tile reversal.



\*What could be the problem of each heuristic?

\*What happens if a heuristic function returns non-unique and conflicting scores?

\* Devising good heuristics based on the limited information is not easy but critical in solving a complex problem!



#### How can we utilize f(n) in DFS?

Goal

# **Hill-climbing**





# "Landscape" of search for max value

# **Simple Hill-climbing search**



# Algorithm sketch

- 1. Check if the current state is a solution, if so return it else go to next step.
- 2. Expand the current state of the search.
- 3. Evaluate its children states using a heuristic function.
- 4. Select the <u>FIRST</u> better state for further expansion. Continue steps 1 - 4 until it finds a solution or reach no better state.

# **Simple Hill-climbing** as a Heuristic Search

## Local search

- Keep track of single current state
- Move only to neighboring states
- Ignore paths

## Advantages:

- Use very little memory
- Can often find reasonable solutions in large or infinite (continuous) state spaces.
- One of the simplest search methods based on heuristics. It works like DFS that utilizes a heuristic.



#### Local maxima

may get stuck and may fail to find the best solution when it reaches a state that has no other better states.

#### Plateau

may get confused by the result of evaluation when the best is not clear

#### Shoulder

may get confused by the result of evaluation when the best is not the direct successor

18

# **Variations of Hill-climbing**



## • Simple Hill-climbing focusing on step 3

- 1. Check if the current state is a solution, if so return it else go to next step.
- 2. Expand the current state of the search.
- 3. Evaluate its children states.
- 4. Select the <u>FIRST</u> <u>better</u> state for further expansion.

Continue steps 1 – 4 until it finds a solution or reach a no better state.

### • Improvement at step 3

- All successors are compared and select the <u>best state</u>.
  - Called <u>Steepest-ascent/descent</u> Hill-climbing or Gradient Search
    - The decision on descent or ascent **depends on the heuristics**.

# Gradient ascent/descent search

function HILL-CLIMBING( problem) return a state that is a local maximum
input: problem, a problem
local variables: current, a node.
neighbor, a node.

*current* ← MAKE-NODE(INITIAL-STATE[*problem*])

loop do

 $\begin{array}{l} \textit{neighbor} \leftarrow a \text{ highest valued successor of } \textit{current} \\ \textbf{if VALUE} [\textit{neighbor}] \leq \text{VALUE}[\textit{current}] \textbf{ then return STATE}[\textit{current}] \\ \textit{current} \leftarrow \textit{neighbor} \end{array}$ 

## Plateau Problem in Hill-Climbing with 3-Level Look Ahead





\*Hill-climbing can get confused in this case (**plateau**) as the cost of all paths are the same or similar.

How can we handle the local maxima problem?

# Student Participation: Hillclimbing



Can these variations of Hill-climbing solve local maximum, shoulder, plateau problems?

# Simulated Annealing (SA)



- The name and inspiration of SA come from annealing in metallurgy, a technique involving heating and controlled cooling of a material to reduce their defects.
- A generic <u>probabilistic</u> meta-heuristic for the **global optimization problem**.



# Physical Interpretation of Simulated Annealing



- Annealing = physical process of cooling a liquid or metal until particles achieve a certain frozen crystal state
  - simulated annealing:
    - free variables are like particles
    - seek "low energy" (high quality) configuration
    - get this by slowly reducing temperature T, which particles move around randomly

# Is SA search able to handle the local maxima?



Despite the many local maxima in this graph, the global maximum can still be found using **simulated annealing**.

# **Search using Simulated Annealing**

- Simulated Annealing = hill-climbing with non-deterministic search
- Basic ideas:
  - like hill-climbing identify the quality of the local improvements
  - instead of picking the best move, pick one randomly
  - say the change in objective function is  $\Delta$  (big delta)
  - if  $\Delta$  is positive, then move to that state
  - otherwise:
    - move to this state with **probability proportional to**  $\Delta$
    - thus: worse moves (very large negative  $\Delta$ ) are executed less often
  - however, there is always a chance of escaping from local maxima over time, make it less likely to accept locally bad moves
  - Can also make the size of the move random as well, i.e., allow "large" steps in state space

# **Simulated Annealing**

function SIMULATED-ANNEALING( problem, schedule) return a solution state

input: problem, a problem

schedule, a mapping from time to temperature

local variables: current, a node.

next, a node.

*T*, a "temperature" controlling the probability of downward steps

```
current ← MAKE-NODE(INITIAL-STATE[problem])
```

for  $t \leftarrow 1$  to  $\infty$  do

 $T \leftarrow schedule[t]$ 

if T = 0 then return *current* 

 $\mathit{next} \leftarrow a randomly selected successor of current$ 

 $\Delta E \leftarrow VALUE[next] - VALUE[current]$ 

if  $\Delta E > 0$  then current  $\leftarrow$  next

else current  $\leftarrow$  next only with probability  $e^{\Delta E/T}$ 

### **Algorithm sketch**

At each step, the SA heuristic considers some neighbor, s' of the current state s, and *probabilistically* decides moving from s to s'.

The probabilities are chosen so that the system ultimately tends to move to states of **lower energy** (annealing schedule,  $p = e^{\Delta E/T}$ ).

Defining an annealing schedule, p is critical.

This step is repeated until the system reaches a state that is solution, good enough for the application, or until a given computation budget has been exhausted.



# **More Details on Simulated Annealing**

- Lets say there are 3 moves available, with changes in the | objective function of  $\Delta E_1 = -0.1$ ,  $\Delta E_2 = 0.5$ ,  $\Delta E_3 = -5$ . (Let T = 1).
- pick a move randomly:
  - if  $\Delta E_2$  is picked, move there.
  - if  $\Delta E_1$  or  $\Delta E_3$  are picked, probability of move = exp( $\Delta E/T$ )
  - move 1: prob1 = exp(-0.1) = 0.9,
    - i.e., 90% of the time we will accept this move
  - move 3: prob3 = exp(-5) = 0.05
    - i.e., 5% of the time we will accept this move
- T = "temperature" parameter
  - high T => probability of "locally bad" move is higher
  - low T => probability of "locally bad" move is lower
  - typically, T is decreased as the algorithm runs longer
    - i.e., there is a "temperature schedule"

# **Simulated Annealing in Practice**



- The method was proposed in 1983 by IBM researchers for solving VLSI layout problems (Kirkpatrick et al, Science, 220:671-680, 1983).
  - theoretically will always find the global optimum (the best solution)

# **Problems of SA**



- Useful for some problems, but can be very **slow** 
  - slowness comes about because T must be decreased very gradually to retain optimality
  - In practice how do we decide the rate at which to decrease T? (this is a practical problem with this method)
- Unfortunately, the applicability of simulated annealing is problem-specific because it relies on finding lucky jumps that improve the position. In such extreme examples, hill climbing will most probably produce a local maxima.

# References



- George Fluger, Artificial Intelligence: Structures and Strategies for Complex Problem Solving, 6<sup>th</sup> edition, Chapter 4, Addison Wesley, 2009.
- Russel and Norvig, Artificial Intelligence: A Modern Approach, 3<sup>rd</sup> edition, Prentice Hall, 2010.